# Nonrandomness of the 33-round MD6 

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## Differentiability from the random oracle

Fix $n$ bits of input and $n$ bits of output.


- To find a solution for a random oracle one needs $2^{n}$ trials;
- For some compression functions we can generate solutions with cost 1.


## MD6 compression function

- H: 89 words $\rightarrow 16$ words.
- Write down bit equations for the compression function:

■ Nonlinear: $x_{i-89}^{j}+x_{i-31}^{j} x_{i-67}^{j}+x_{i-18}^{j} x_{i-21}^{j}+x_{i-17}^{j}+x_{i}^{j}=0$.

- $2 \times$ Linear: $y_{i}^{j}=x_{i}^{j}+x_{i}^{j+l_{i}}$.
- $1664 \times 3 \times 64 \approx 300000$ equations for MD6-256.
- Set some variables to constants.
- Solve a system with a Gaussian-like process.
- Generate many solutions.


## Results

We fix several bits in the input and the output of the compression function - and show how to derive the others.

| Rounds | Fixed bits |  | Speed |  |
| :---: | :---: | :---: | :---: | :---: |
|  | input | output | 32-bit | 64-bit |
| 18 | Aumasson et al. |  | 11 | 5 |
| 22 | > 9 | $>9$ | 14 | 6 |
| 26 | 6 | 6 | 17 | 7 |
| 30 | 2 | 2 | 20 | 8 |
| 32 | 2 | 2 | 21 | 9 |
| 33 | 1 | 1 | 22 | 9 |
| 80 | MD6-160 |  | 52 | 22 |
| 96 | MD6-224 |  | 63 | 26 |
| 104 | MD6-256 |  | 68 | 28 |

## Tool


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