



“Whitening 2 Last 32 Bit Hash Messages

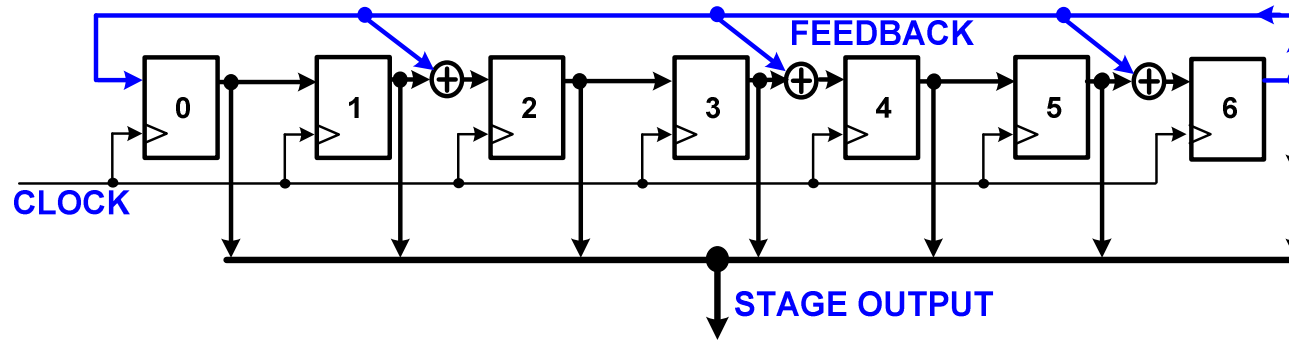
with a HAIFA* Inspired 64 Bit Hybrid Mersenne

Prime Number LFSRs/Binary Counter”

An Efficient ZK-Crypt Artifact

* **Eli Biham & Orr Dunkelman, “A Framework for Iterative Hash Functions–HAIFA”, Technion, 2006**
Patents Pending

A 7 Celled Mersenne Prime Number LFSR



MERSENNE LFSR - 7 CELL - 127 UNIQUE PSEUDO RANDOM OUTPUT STAGES

ONE TO MANY CONFIG WITH TAPS 1, 3, 5, & 6- MAX DISTANCE BETWEEN TAPS = 2

AS CLOCK PULSE RISES-INPUT SHIFTS TO OUTPUT OF EACH (FLIP-FLOP) CELL.

INITIAL CONDITION- ALL CELLS ARE SET TO '1'.

A 7 BIT BINARY COUNTER'S GATE COUNT IS 87 - THE 7 BIT LFSR NEEDS 51 GATES
→ $36/87 = 41\%$ FEWER GATES

A 64 BIT BINARY UP COUNTER HAS MAX PROPAGATION TIME FOREVER.

12/02/09 12:12
1 A 7 Bit One to Many Mersenne Prime LFSR.vsd



Why 1 to Many Galois LFSRs for Unique Counting

LFSRs are Efficient Large Number Counters
40% Fewer Gates than Same Size Binary Counters
Faster No Delays - No Ripple, No Carry
Almost No Bias on any Bit - $2^n/2$ '1's $2^n/2-1$ '0's
Each LFSR has 2^n-1 Unique Pseudo Random Stages

1 to Many LFSRs are "Whiter" than Many to '1's
Less Correlated Motion Sense than Many to '1's
More Local Pseudo-Randomness Best if Taps are
not Overly Distanced from Nearest Neighbor



What About Mersenne Prime Number LFSRs

Mersenne LFSRs Have a Prime Number of Stages

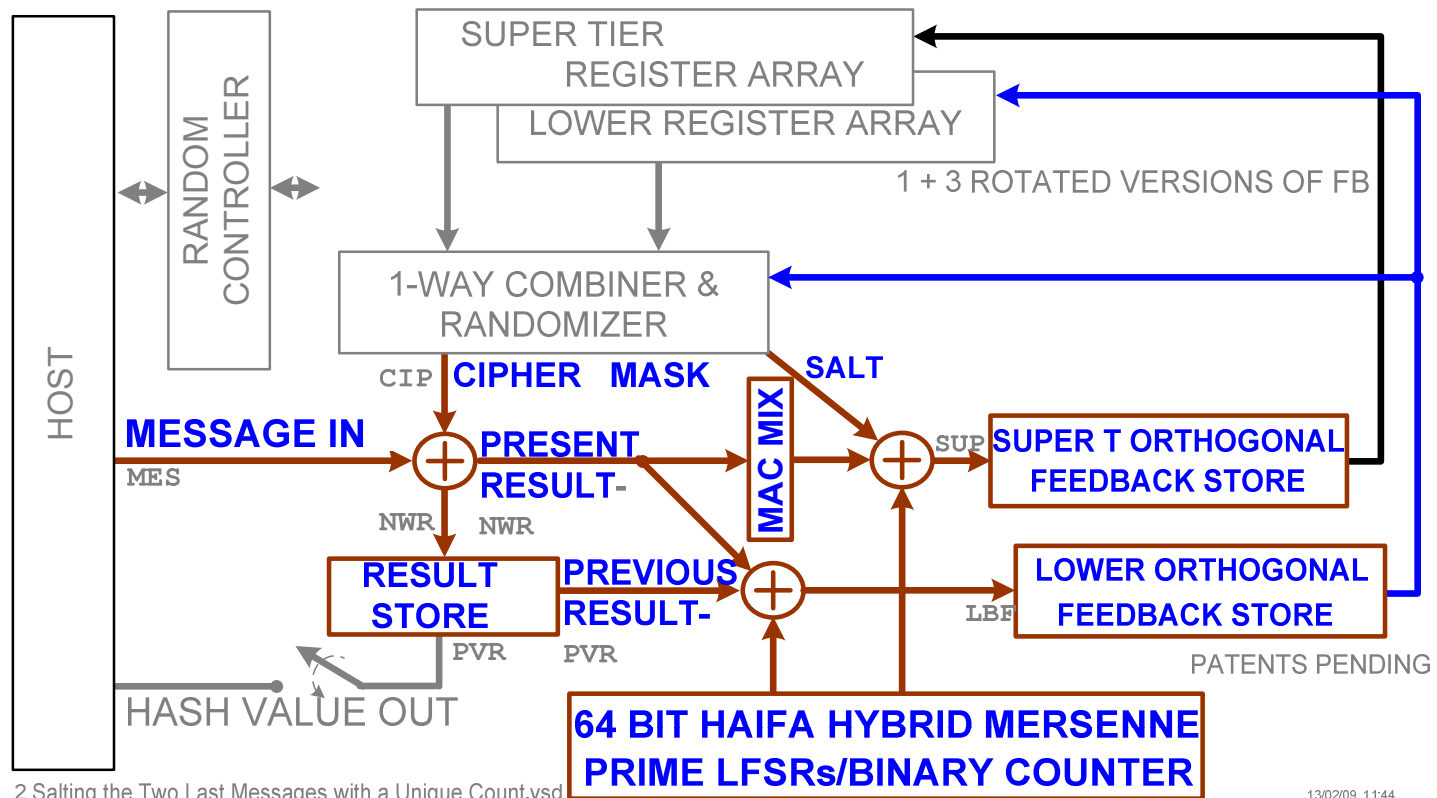
CoPrime to Each Other & Relatively Prime LFSRs
can be Concatenated to one Even Length Counter

All Can be Concatenated into One Large Counter

The Few Mersennes 2,3,5,7,13,17,19,31 & 61(?)

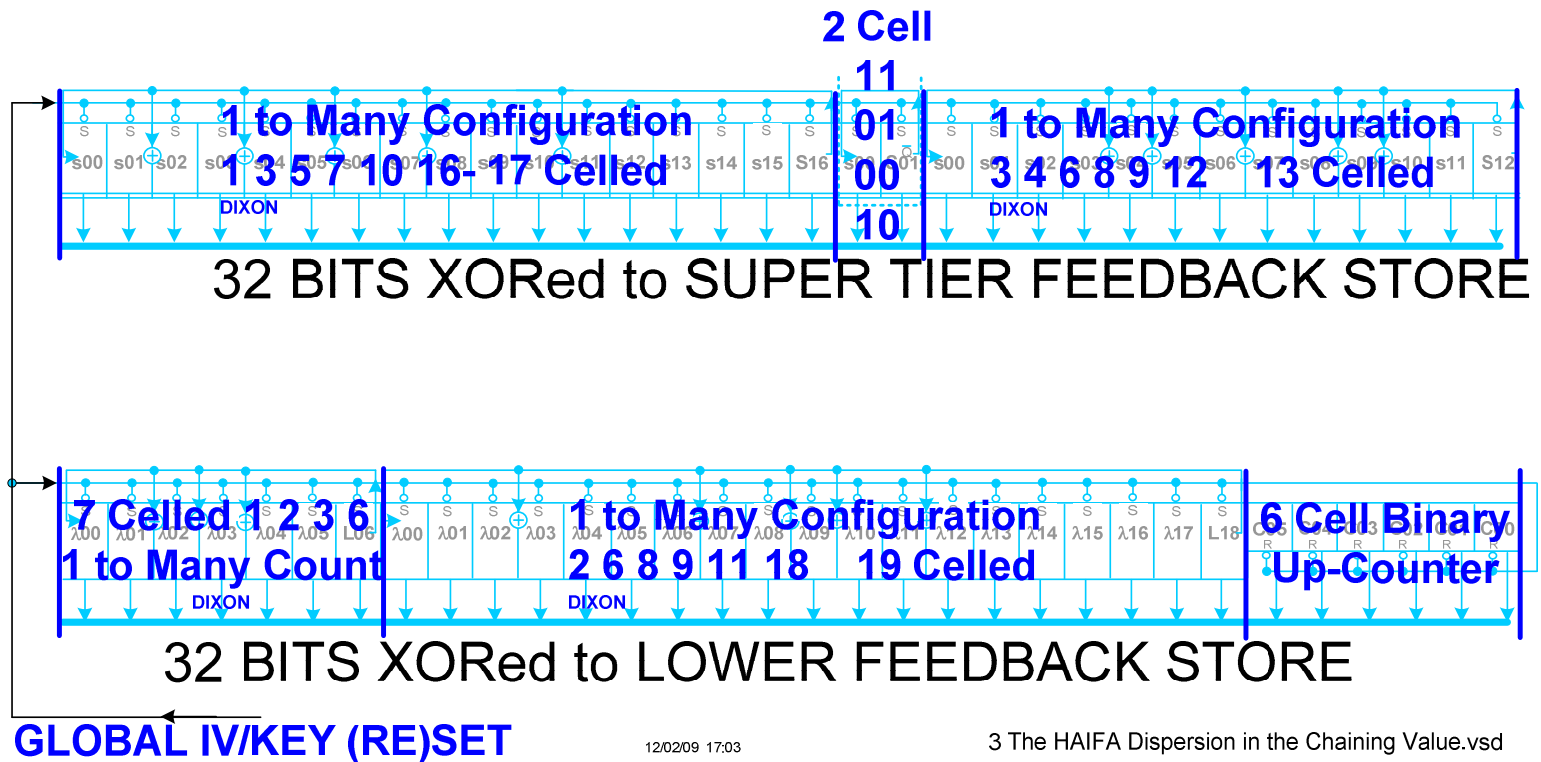
<7 Cells too Biased to '1'; 31 Bit Cells too few Taps

Mersenne was the Father of the Math of Music- This Looks Like Ultra Modern Symmetric Dissonance



M COUNTERS MAKE 64 BIT FIXED POINTS IN CHAINING VALUES

Did Eli or Orr Anticipate a 64 Bit Balanced Count





For Expansion PRFs Merkle-Damgård Loses Entropy

Unique Message Counts - from -

Super Tier FB - $131,071 \times 8191 = 1,073,602,561$

Lwr FB - $127 \times 524,287 \times 64 = 4,261,404,736$

Unique Stages in Multiple of 2 Counters =
 4.58×10^{18}

$\sim 2^{67}$ Processed Data Bits

A 2^{62} Binary Counter = 4.61×10^{18} no big loss.



You Couldn't Fall Asleep in 5 Minutes So-

Thanks for Your (Prime- Indivisible) Attention

Thanks to Relative Prime Counters

Thanks to Eli, Orr and Hugo for Inspiration

The ZK-Crypt Design Group - FortressGB

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